

# CSE 332 Final

Olcay Taner YILDIZ

## I. QUESTION (16 POINTS)

Consider the following snapshot of a system, which has four types of resources A, B, C and D and five processes  $P_1$ ,  $P_2$ ,  $P_3$ ,  $P_4$  and  $P_5$ .

	Allocation	Max	Available
	-----	-----	-----
	A B C D	A B C D	A B C D
P1	1 0 0 0	2 6 5 0	1 0 2 2
P2	1 3 4 4	2 3 7 6	
P3	0 2 3 2	0 4 5 2	
P4	0 1 1 4	0 4 5 6	
P5	0 0 1 1	1 0 1 2	

Answer the following questions using the Bankers algorithm:

- Is the system in a safe state? If so, provide a safe sequence. If not, justify your answer.
- If a request from process P1 arrives for (1,0,1,0), should the request be immediately granted?

## II. QUESTION (12 POINTS)

Given two processes, ( $P_1$ ,  $P_2$ ) and four statements ( $S_1$ ,  $S_2$ ,  $S_3$ ,  $S_4$ ), where  $P_1$  will execute  $S_1$  and  $S_2$ , and  $P_2$  will execute  $S_3$  and  $S_4$ . Use semaphores to make the order of execution  $S_2$ ,  $S_3$ ,  $S_1$ ,  $S_4$ . Use at most four semaphores. Show the initial semaphore values.

## III. QUESTION (15 POINTS)

The first known correct software solution to the critical-section problem for two processes was developed by Dekker. The two processes,  $P_0$  and  $P_1$ , share the following variables, boolean flag[2] and int turn. The structure of process  $P_i$  ( $i == 0$  or  $1$ ) is shown below; the other process is  $P_j$  ( $j == 1$  or  $0$ ). Prove that the algorithm satisfies all three requirements for the critical section problem.

```
do {
    flag[i] = TRUE;
    while (flag[j]) {
        if (turn == j) {
            flag[i] = false;
            while (turn == j)
                ; // do nothing
            flag[i] = TRUE;
        }
    }
    // critical section
    turn = j;
    flag[i] = FALSE;
    // remainder section
}while (TRUE);
```

## IV. QUESTION (14 POINTS)

Draw the resource allocation graph for the following situation. There are three processes  $P_1$ ,  $P_2$ , and  $P_3$ . There are four resources  $R_1$ ,  $R_2$ ,  $R_3$ , and  $R_4$ . Resources  $R_2$ ,  $R_1$ ,  $R_4$  and  $R_3$  are allocated to processes  $P_1$ ,  $P_2$ ,  $P_2$  and  $P_3$  respectively. Processes  $P_1$  and  $P_2$  request resources  $R_1$  and  $R_3$  respectively. Is there a deadlock? If not, suggest a single change in the system, so that the resulting system is in deadlock state.

## V. QUESTION (12 POINTS)

- A computer with a 32bit address uses a twolevel page table. Virtual addresses are split into a 9bit toplevel page table field, and 11bit secondlevel page table field, and an offset. How large are the pages? How many pages are there in the virtual address space?
- A computer whose processes have 1024 pages in their address space keeps its page tables in memory. The overhead required for reading a word from the page table is 500 nsec. To reduce this overhead, the computer has an associative memory, which holds 32 (virtual page, physical page frame) pairs, and can do a look up in 100 nsec. We say that there is a cache hit if a page table entry needed to handle a reference to a virtual page is in the associative memory. What hit rate is needed to reduce the mean overhead to 200 nsec?

## VI. QUESTION (15 POINTS)

Consider the following page reference string:

1, 2, 3, 4, 2, 5, 7, 2, 1, 3, 6, 2, 2, 3, 2

How many page faults occur for the replacement algorithms LRU, FIFO and Optimal, assuming there are 3 page frames? Assume that all the frames are initially empty.

## VII. QUESTION (12 POINTS)

Consider the two-dimensional array A:

```
int A[][] = new int[100][100];
```

where  $A[0][0]$  is at location 200 in a paged memory system with pages of size 200. For three page frames, how many page faults are generated by the following array-initialization loops?

- for ( $j = 0; j < 100; j++$ )  
 for ( $i = 0; i < 100; i++$ )  
 $A[i][j] = 0;$
- for ( $i = 0; i < 100; i++$ )  
 for ( $j = 0; j < 100; j++$ )  
 $A[i][j] = 0;$